Controlled conflict resolution for replicated document

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Collaboration on Shared document

- Mutable shared data in distributed system
- CAP theorem [Brewer 2000]
 - strong Consistency
 - Availability
 - High responsiveness
 - Partition tolerance
 - Mobile Disconnected work
- A solution : Optimistic Replication (Eventual consistency)

Optimistic replication

- A replica per application instance
- Immediately available for updates
- Updates sent to other replicas
- Consistency Model
 - (Strong) Eventual Consistency
- Replicated merge mechanism
 - Concurrent modifications
 - Possibly "conflicting"

Concurrency Conflicts

Replica A

• Removes an element

Adds a paragraph P

Replica B

• Removes the element and readds it (e.g. undo)

• Inserts a text • Inserts another text at the more complexity position

more conflicts

moves the section P

Sets a title • Sets a title

Proof on ad-hoc complex systems

- Commutative Replicated Data Type $\forall op_1, op_2 \in Op, S \text{ is state}$ $op_1 || op_2 \Rightarrow [op_1, op_2](s) = [op_2, op_1](s)$
- Operational Transformation $\forall op_1, op_2 \in Op, S$ is state $op_1 || op_2 \Rightarrow [op_1, IT(op_2, op_1)](s) = [op_2, IT(op_1, op_2)](s)$ $op \in Op,$ $IT(IT(op, op_1), IT(op_2, op_1)) = IT(IT(op, op_2), IT(op_1, op_2))$
 - More different operation \rightarrow more complex proof

Our proposition



- Decouples eventual consistency from conflict resolution
- Constructs a new data type using simpler one
- Layered approach
 - Bottom layers for eventual consistency
 - Layers per data type constraint

Conflicts to constraints

Replica A

- Removes an element
- Inserts a text

Sets a title

Replica B

• Removes the element and readds it (e.g. undo)

It's a set

• Inserts another text at the same position

Elements are totally ordered

• Adds a paragraph *P*

• Removes the section *P* belongs

. . .

• Sets a title

All elements are connected

Only one title

Layered approach Example : file system data



Stateless Layers



- Replication layer in charge of merge
- Adaptation layer computes a lookup
- Straight forward strong eventual consistency
- Fits state-based replication

Incremental Layers



- Adaptation layer state is updated
- No concurrency management in adaptation layer
 "all local"
- Fits operation-based replication
- Proof is more difficult. (proof that it's equivalent of stateless layers)

Adaptation layer example: Trees policies



Incremental reappear policy

- Mark reappeared nodes "ghosts"
- Path added in the inner set
 - Creates *recursively* missing father(s)
 - Creates node if needed (else unmark)
- Path removed in the inner set
 - If corresponding node got no child
 - Removes it
 - If it was the last child of a reappeared node, removes *recursively* its father
 - Else marks it



Example of incremental reappear policy



- ➢ Del(ab)
- Del(abc)
- ➤ Del(abd)
- > Add(abc)

Ghost without children is deleted

Ordered Trees



- Position path are list of couple (position, label)
- Position are unique, totally ordered and define in a dense space
- Several existing possibility
 - Logoot [Weiss et al 09],
 - FCEdit [Martin, Lugiez 09],
 - Treedoc [Preguiça et al 09],
 - PPS [Wu et al 10], ...

Experiments

- Evaluate performance of the approach
 - Overhead of layers and connection algorithms
 - 4 connection policies : skip, reappear, root, compact
 - 2 positions identifier : Logoot, WOOTH
- Skip policy compared to
 - Operational Transformation (OT) [Ressel et al 96]
 - "Naïve" tree based on TTF [Oster et al 06]
 - TreeOPT [Ignat, Morrie 03]
 - Conflict-free Replicated Data Types (CRDT) [Shapiro et al 11]
 - Tree CRDT : FCEdit [Martin, Lugiez 09]

Experiments Set up

• Open-source Benchmark (Java)

http://github.com/score-team/replication-benchmarker

• Layered implementation

new PositionIdentifierTree(new WordTree(new ReappearPolicy(), new CounterSet()))

- One randomly generated trace (available)
 - 30,000 operations
 - Random paths
 - 80% insertions
- Intel(R) Xeon(R) 5160 dual-core
 - 3 executions

Experiments : Local execution time skip policy 35 FCEdit OTTree --Logoot WOOTH 30 TreeOPT -25 Time (in microseconds) 20 15 10 5 0 50 100 150 200 250 0 300 Operation (100x) All approaches are below 30 μ s •

Experiments : Remote execution time Skip policy (log scale)



Experiments : Memory Skip policy (log scale)







• All policies (except root) are below 120 μ s

Experiments : Remote execution times policies



Related Work

- Replicated data types
 - Ad-hoc, OT, CRDT, ...
 - We use them as replication layer
- Generic replication frameworks
 - Bayou [Terry et al 95], IceCube [Kermarrec et al 01]
 - Do not scale well since centralized
- Collaborative applications
 - DVCS (git, darcs, hg, ...), replicated file systems (Fuse,
 - Some conflicts have to be resolved by the users



Conclusion



- Benefit of our approach
 - Behavior modularity
 - More complex data types (XML, ...)
 - "Conflict outside the data" : separate awareness from consistency
- Performances
 - Very stable
 - Comparable to ad-hoc approaches

Future Work



- Experiments on large-scale real data
 - Open-sources DVCS traces
- User file-system (FUSE)
- Incremental DTD correction layers
- Effective awareness support
- Formal proof of equivalence between incremental and stateless policies